

A TOPOS-THEORETIC APPROACH TO PARACONSISTENT POSSIBILISTIC LOGIC

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- 2 POSSIBILISTIC LOGIC
- 3 TOPOS THEORY
- 4 PUTTING IT ALL TOGETHER
- 5 STILL OTHER APPROACHES

LOGIC AT CLE

- Research developed under supervision of Walter Carnielli.

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Plurality of systems: **C**-systems, LFI's, etc.

ALGEBRA AND TOPOLOGY

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Kind of duality: closed set topology stands for paraconsistent logic as open set topology stands for intuitionistic logic.

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Possibilistic logic expression (ϕ, α) is understood as $N(\phi) \geq \alpha$.

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Given $K = \{(\phi_i, \alpha_i)\}_{i=1, \dots, n}$, one defines level of inconsistency

$Inc(K) = \max\{\alpha \mid K \vdash (\perp, \alpha)\}$.

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Very interesting and useful feature: comes with logic!

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A *complement classifier* for a category \mathcal{E} with terminal object 1 is an object Ω together with an arrow $F : 1 \rightarrow \Omega$ satisfying the condition that for every monic arrow $f : a \rightarrow b$ there exists a unique arrow $\bar{\chi}_f$ such that the following diagram is a pullback

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$\bar{\chi}_f$ is the complement character of f . \mathcal{E} is a topos whose truth-values form a paraconsistent algebra.

TOPOI AND POSSIBILISTIC LOGIC - 1

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For $L \in \mathcal{B}$, we define the *possibility of L* (according to π) as

$$\Pi_{\pi}(L) = \bigvee_{\omega \models L} \pi(\omega).$$

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 $\mathbf{n} = \mathbf{0}$: for every $L \in \mathcal{B}$ and $r \in [0, 1]$, the following are objects of \mathbf{U} : $\{\pi : \Pi_\pi(L) = r\}$, $\{\pi : \pi \leq \mathbf{r}\}$

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$\mathbf{n} + \mathbf{1}$: if A, B and A_i are objects of \mathbf{U} then the following are

objects of \mathbf{U} : $\{1 - \pi : \pi \in A\}, \cap A_i, \{\pi_1 \times \pi_2 : \pi_1 \in A, \pi_2 \in B\},$
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THEOREM (C. SOSSAI)

Presheaves over \mathbf{U} can represent possibilistic logic.

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Introduce a topology in \mathbf{U} , and then use closed set sheaves.

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Modifying Sossai's approach: at least three different choices:
Construct complement classifier in \mathbf{U} -Sets, and investigate the obtained logic.

Introduce a topology in \mathbf{U} , and then use closed set sheaves.

Use categorical topologies on \mathbf{U} or \mathbf{U} -sets.

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Can we do something similar for possibility theory?

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Thank you!
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