

Recall: Let  $M$  be a rank- $r$  matroid.

- All circuits have size at most  $r+1$
- $M$  is uniform if all circuits have size  $r+1$
- $M$  is  paving if all circuits have size  $r$  or  $r+1$ .

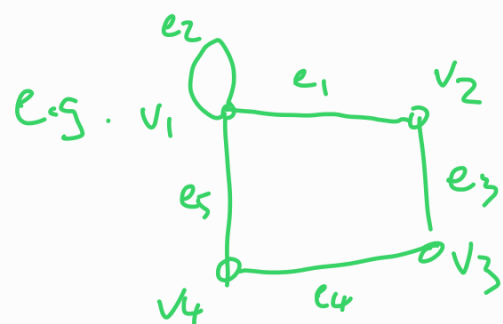
Recap: last time we saw matroids that arise from a set of vectors in  $\mathbb{F}^m$  (or, a matrix over  $\mathbb{F}$ ) - these are called  $\mathbb{F}$ -representable matroids

Today, matroids that arise from graphs.

Recall: a graph consists of  $G = (V, E)$

- a set  $V$  called the vertex set
- a set  $E$  called the edge set

and an incidence function  $\varphi$  that maps each  $e \in E$  to either a pair in  $V$ , or a singleton in  $V$



$$V = \{v_1, v_2, v_3, v_4\}$$

$$E = \{e_1, \dots, e_5\}$$

$$\varphi(e_1) = \{v_1, v_2\}, \quad \varphi(e_2) = \{v_1\},$$

...

adjacent edges

when a vertex and edge are incident

ends of edge

a walk in a graph

path

closed walk

- first and last vertex are the same.

cycle

a subgraph of  $G = (V, E)$  is a graph  $G' = (V', E')$   
where  $V' \subseteq V$  and  $E' \subseteq E$   
and  $\mathcal{Q}(e') \subseteq V'$  for all  $e' \in E'$ .

a subgraph  $G'$  of  $G = (V, E)$  is edge-induced,  
for some  $E' \subseteq E$  when  $G' = (V', E')$  with

$V' = \bigcup_{e \in E'} \mathcal{Q}(e)$ . We denote  $G'$  as  $G[E']$ .

Theorem: Let  $G$  be a graph with edge set  $E$ .

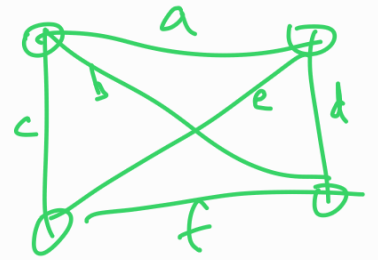
Let  $\mathcal{C}$  be

$\{C \subseteq E : C \text{ is the edge set of a cycle in } G\}$ .

Then  $\mathcal{L}$  is the family of circuits of a matroid on ground set  $E$ .

We denote this matroid as  $M(G)$ .

e.g. Consider the graph  $K_4$  on edge set  $E = \{a, b, c, d, e, f\}$  as illustrated.

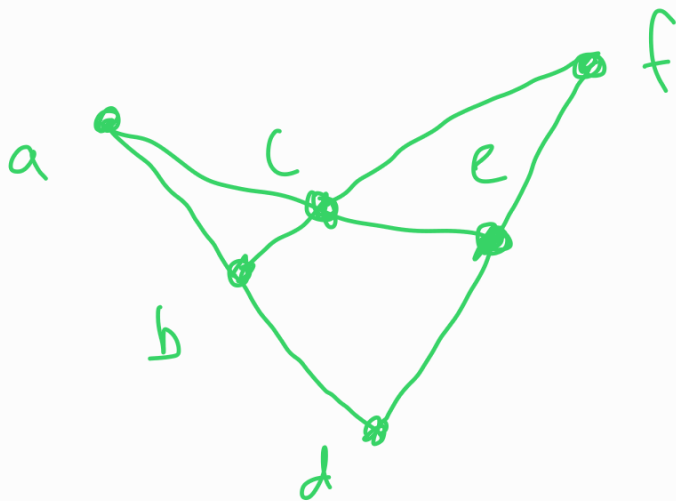


$$\mathcal{L} = \{ \{a, b, d\}, \{a, c, e\}, \{d, e, f\}, \{b, c, f\}, \\ \{a, c, d, f\}, \{a, b, e, f\}, \{b, c, d, e\} \}$$

The bases have size 3, so

$M(K_4)$  has rank 3.

A geometric representation for  $M(K_4)$ :

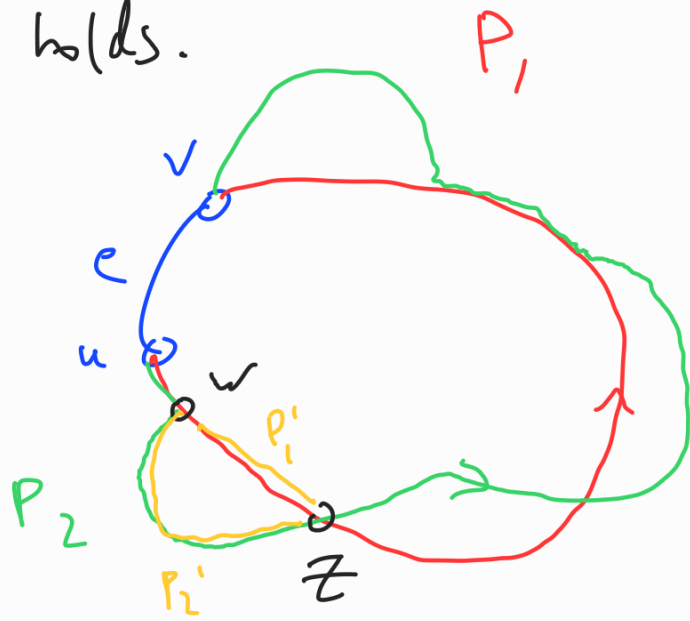


Proof of  $\mathcal{P}^n$ : We will show  $\mathcal{C}$  satisfies (C1)-(C3).

Since a cycle of a graph contains at least one edge,  $\emptyset \notin \mathcal{C}$ , so (C1) holds. Moreover, no cycle in a graph contains a smaller cycle as a proper subgraph, so (C2) holds.

It remains to show (C3) holds.

Let  $C_1$  and  $C_2$  be edge sets of distinct cycles in  $\mathcal{C}$ , with  $e \in C_1 \cap C_2$ . We will show that there is a cycle on edge set  $C_3$  where  $C_3 \subseteq (C_1 \cup C_2) - \{e\}$ .



Note that  $e$  is not a loop, otherwise  $C_1 = C_2 = \{e\}$ . Let  $u$  and  $v$  be the ends of  $e$ . Let  $P_i$  be the path from  $u$  to  $v$  along the edges in  $C_i - \{e\}$ , for  $i \in \{1, 2\}$ . Let  $P_1 = v_0, e_1, v_1, \dots, v_k$

and  $P_2 = v_0', e_1', v_1', \dots, v_{k_1}'$

where  $v_0 = v_0' = u$  and  $v_k = v_{k_1}' = v$ .

Let  $w$  be the first vertex  $v_i$  in  $P_1$  for which  $e_{i+1}$  is not in  $P_2$ .

Traverse  $P_1$ , starting at  $w$ , towards  $v$ , and let  $z$  be the first vertex distinct from  $w$  that is in  $P_2$  (such a vertex exists since  $v$  is one such vertex).

Call this path  $P_1'$ . Now let  $P_2'$  be the subpath of  $P_2$  obtained by starting at  $w$  and ending at  $z$ .

Now  $P_1'$  and  $P_2'$  are both paths from  $w$  to  $z$ , with no other vertices in common, so letting  $C_3$  be the union of the edge sets of  $P_1'$  and  $P_2'$ ,  $C_3$  is the edge set of a cycle contained in  $(C_1 \cup C_2) - \{e\}$ .

Thus (C3) holds.  $\square$

For a graph  $G$ , we call this matrix  $M(G)$  the cycle matrix of  $G$ .

A matrix  $M$  is graphic if there exists a graph  $G$  such that  $M \cong M(G)$ .

For a graph  $G$ , a forest of  $G$  is a set  $F \subseteq E(G)$  such that  $G[F]$  contains no cycles,

and a forest  $F$  is maximal if there is no proper superset  $F' \subseteq E(G)$  of  $F$  such that  $F'$  is a forest of  $G$ .

For a graph  $G$

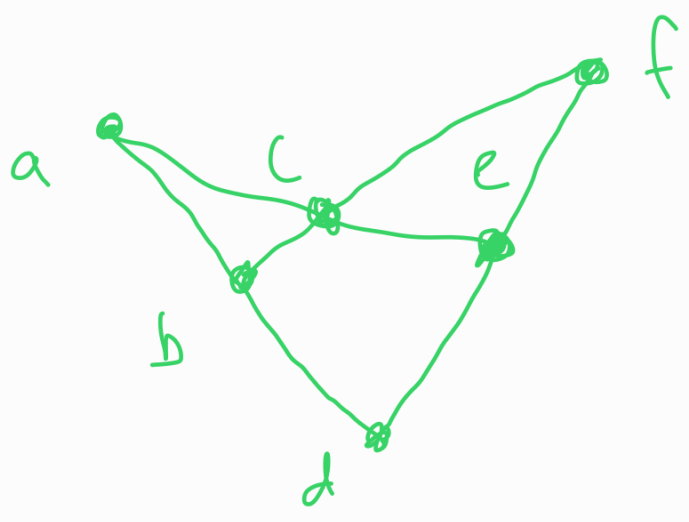
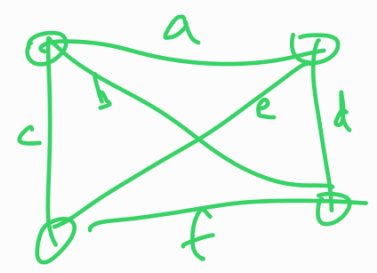
$\left. \begin{array}{l} \text{circuits in } M(G) \\ \text{independent sets} \\ \text{bases} \end{array} \right\}$  correspond to  $\left\{ \begin{array}{l} \text{cycles in } G \\ \text{forests of } G \\ \text{maximal forests} \\ \text{of } G. \end{array} \right.$  respectively.

When  $G$  is connected, the bases of  $M(G)$  are the edge sets of spanning trees of  $G$ .

A spanning tree of a graph  $G$  is a subgraph  $G'$  of  $G$  such that  $G'$  is a tree and  $V(G) = V(G')$

visited.

e.g.<sup>1</sup> Consider the graph  $K_4$  on edge set  $E = \{a, b, c, d, e, f\}$  as illustrated.



$$A = \begin{bmatrix} & a & b & c & d & e & f \\ 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$

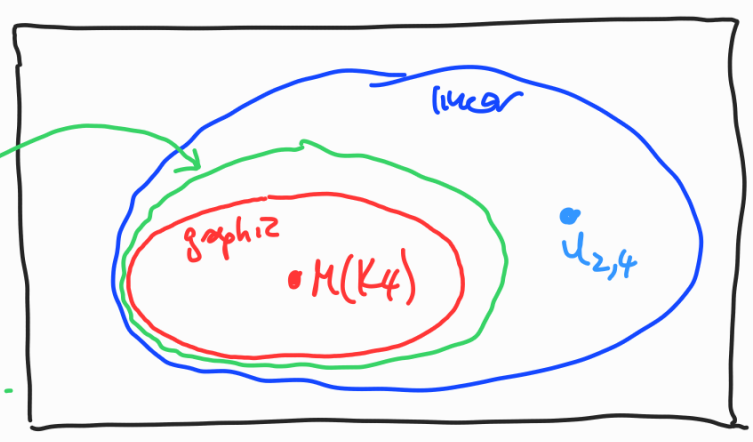
over  $GF(2)$

$$M(K_4) \cong M[A]$$

This demonstrates that  $M(K_4)$  is  $GF(2)$ -representable.

In fact, we will see for any graph  $G$ , the matrix  $M(G)$  is  $GF$ -representable for every field  $GF$ .

regular matrices:  
matrices representable over every field.



Exercise: Prove that  $U_{2,4}$  is not graphic and is not  $\mathbb{F}$ -representable for each field  $\mathbb{F}$ .

However: when  $A = \begin{bmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 2 \end{bmatrix}$  over  $\mathbb{R}$

then  $M[A] \cong U_{2,4}$ , so  $U_{2,4}$  is  $\mathbb{R}$ -representable (and is therefore linear).

Exercise: (1) Characterize for which fields  $\mathbb{F}$  the matroid  $U_{2,4}$  is  $\mathbb{F}$ -representable.

(2) What about  $U_{2,n}$  (for any  $n \geq 2$ ).

Note: in general, characterizing which fields a uniform matroid  $U_{r,n}$  is representable over is a hard problem; no complete answer is known when  $r \geq 5$ .