

Last time: "cryptomorphisms" i.e.

different, but equivalent ways of defining an object - in this case a matroid.

- bases
- independent sets (satisfying (I1) - (I3))
- circuits
- rank
- closure

→ flats

For  $\mathcal{F} \subseteq 2^E$

(F1)  $E \in \mathcal{F}$

(F2) If  $F_1, F_2 \in \mathcal{F}$ , then  $F_1 \cap F_2 \in \mathcal{F}$ .

(F3) If  $F \in \mathcal{F}$  and  $\{F_1, F_2, \dots, F_n\}$  are the minimal members of  $\mathcal{F}$  that properly contain  $F$ , then

$\{F_1 - F, F_2 - F, \dots, F_n - F\}$  is a partition of  $E - F$ .

Theorem 6.5 If  $\mathcal{M}$  is a matroid with family of flats  $\mathcal{F}$ , then  $\mathcal{F}$  satisfies (F1) - (F3). Conversely, if  $E$  is a set and  $\mathcal{F}$  is a family of subsets of  $E$  satisfying (F1) - (F3) then there is a matroid on ground set  $E$  whose family of flats is  $\mathcal{F}$ .

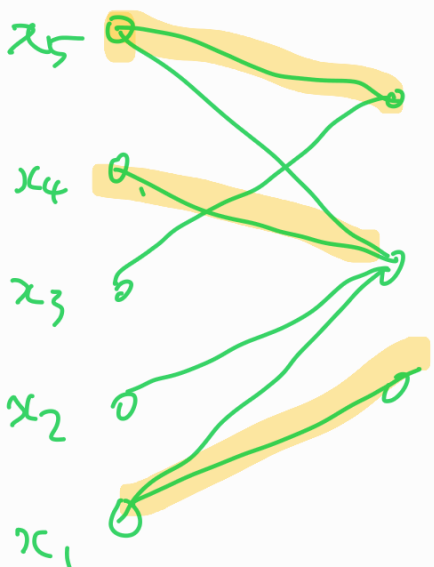
→ independent sets satisfying (I1), (I2), (A1)

For  $\mathcal{I} \subseteq 2^E$ ,

(A1) For all weight functions  $w: E \rightarrow \mathbb{R}$ , the greedy algorithm finds a maximal member of  $\mathcal{I}$  of maximum weight.

Theorem 6.6 Let  $\mathcal{I}$  be a collection of subsets of a set  $E$ . Then  $\mathcal{I}$  is the family of independent sets of a matroid on  $E$  iff  $\mathcal{I}$  satisfies (I1), (I2) and (A1).

Example Consider the transversal matroid on ground set  $\{x_1, x_2, \dots, x_5\}$  with the given presentation and weight function  $w(x_i) = i$  where we seek a solution of maximum weight.



The greedy algorithm considers...

$x_5$  after which  $X = \{x_5\}$ ,

then  $x_4$  after which  $X = \{x_4, x_5\}$ .

then  $x_3$  — we keep  $X = \{x_4, x_5\}$

then  $x_2$  — we keep  $X = \{x_4, x_5\}$

then  $x_1$  after which  $X = \{x_1, x_4, x_5\}$ .

Proof of 6.6 continued from last time.

Last time we saw ( $\Rightarrow$ ). We still need to show

that if  $\mathcal{I}$  satisfies (I1), (I2), (G1), then we have a matroid whose family of independent sets is  $\mathcal{I}$ .

We'll show that if (I3) doesn't hold, then (G1) doesn't hold.

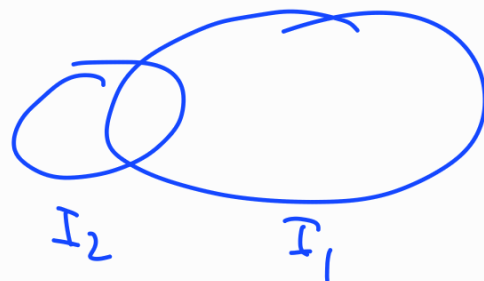
Suppose (I3) fails. Then there exists  $I_1, I_2 \in \mathcal{I}$

with  $|I_2| < |I_1|$  such that  $I_2 \cup e \notin \mathcal{I}$  for all  $e \in I_1 - I_2$ .

Now  $0 \leq |I_2 - I_1| < |I_1 - I_2|$

so we pick  $\varepsilon \in \mathbb{R}$  such that

$$\frac{|I_2 - I_1|}{|I_1 - I_2|} < \varepsilon < 1$$



Now we define  $w: E \rightarrow \mathbb{R}$  as follows:

$$w(e) = \begin{cases} 1 & e \in I_2 \\ \varepsilon & e \in I_1 - I_2 \\ 0 & e \notin I_1 \cup I_2 \end{cases}$$

The greedy algorithm will first pick all elements in  $I_2$ .

Then it considers the elements in  $I_1 - I_2$ , but can't add any of these elements to the partial solution  $X$ , so it will choose a maximal member  $B_G$  of  $\mathcal{I}$  of weight  $|I_2|$ .

But, by (I2),  $I_1$  is contained in a maximal member

$I_1' \in \mathcal{I}$ , and

$$\begin{aligned} w(I_1') &\geq w(I_1) = |I_1 \cap I_2| + \varepsilon |I_1 - I_2| \\ &> |I_1 \cap I_2| + \frac{|I_2 - I_1|}{|I_1 - I_2|} |I_1 - I_2| \\ &= |I_1 \cap I_2| + |I_2 - I_1| \\ &= |I_2| = w(B_G) \end{aligned}$$

Thus, the greedy algorithm fails for this weight function, i.e. (G1) doesn't hold.  $\square$

Representability

Recall:

\* a matroid  $M$  is  $F$ -representable when

$M \cong M[A]$  for some matrix  $A$  over  $F$ .

\*  $M[A] = M[A']$  when  $A'$  is obtained from  $A$  by standard row operations, scaling/swapping columns, applying automorphisms of  $\mathbb{F}$ .

Constructing a representation Let  $M$  be a matroid.

Suppose  $M$  is  $\mathbb{F}$ -representable. Then we know  $M$  has a  $\mathbb{F}$ -representation  $\left[ \begin{array}{c|c} I_r & D \end{array} \right]$  in standard

form where  $B$  labels the columns of  $I_r$ , for some basis  $B$  of  $M$ .

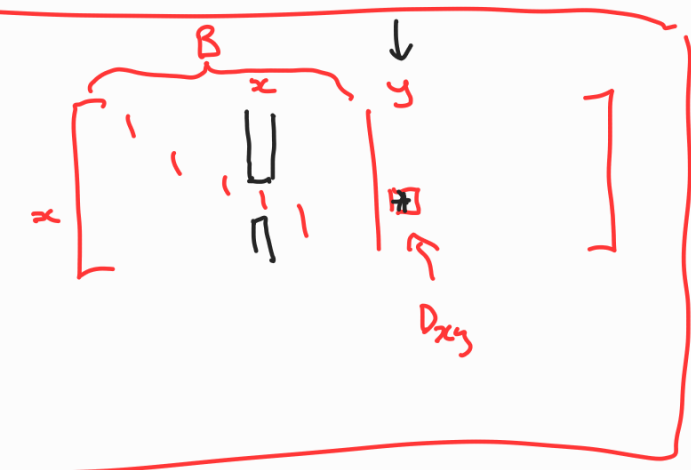
Recall for  $y \in E(M) - B$  that  $C(y, B)$  is the fundamental circuit of  $y$  with respect to  $B$ .

Proposition 7.1: Let  $x \in B$  and  $y \in E(M) - B$

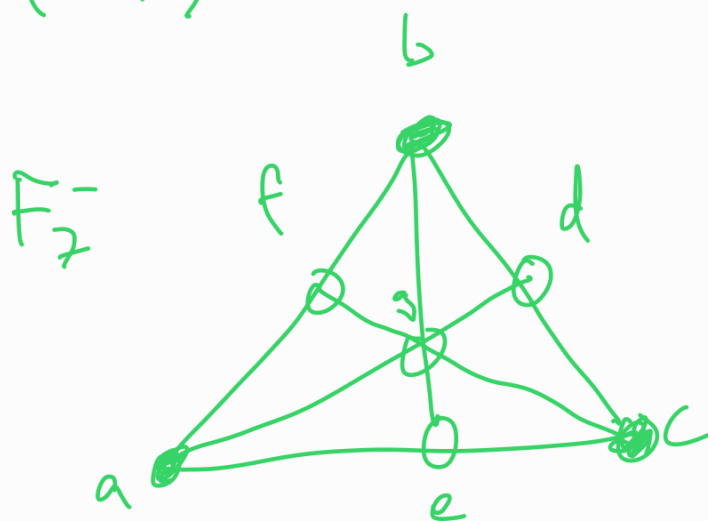
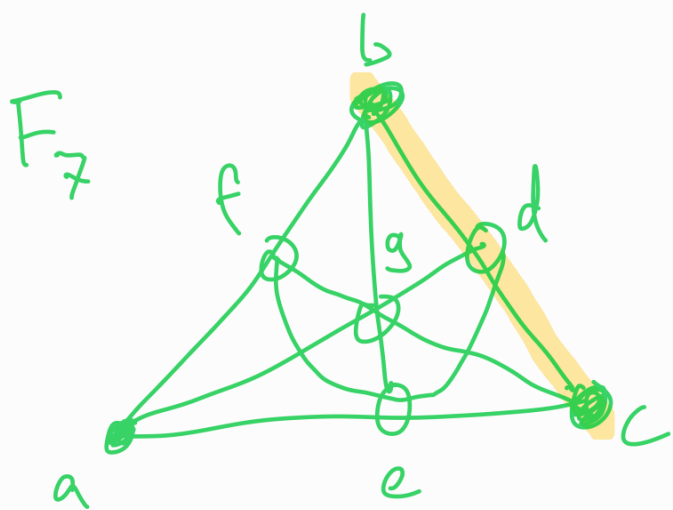
$D_{xy} \neq 0$  if and only if  $x \in C(y, B)$ .

Proof:  $D_{xy} \neq 0 \iff \det(D[\{x\}, \{y\}]) \neq 0$

Prop 3.13  $\Leftrightarrow (B-x) \cup y$  is a basis  
 $\Leftrightarrow (B-x) \cup y$  doesn't contain any circuit  
 $\Leftrightarrow x \in C(y, B)$ . □



Example Consider the Fano matroid  $F_7$ , and the non-Fano matroid  $F_7^-$ .



What fields (if any) are these matroids representable over?

Let  $M \in \{F_7, F_7^-\}$

Suppose  $M$  is GF-representable. Then

it has a  $\mathbb{F}$ -representation of the form:

$$\left[ \begin{array}{ccc|ccc} a & b & c & d & e & f & g \\ 1 & 0 & 0 & 0 & & & \\ 0 & 1 & 0 & * & D & & \\ 0 & 0 & 1 & * & & & \end{array} \right]$$

Consider fundamental circuits with respect to the basis  $B = \{a, b, c\}$ ; these tell us whether entries in  $D$  are zero or non-zero.

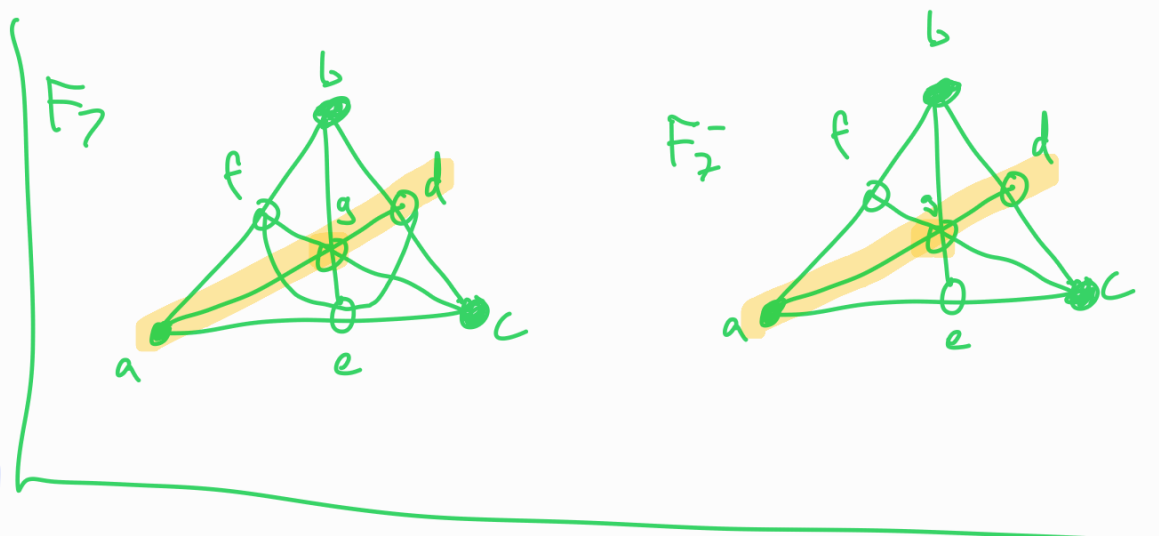
For example, since  $\{b, c, d\}$  is a circuit, the column labelled  $d$  is of the form  $\begin{pmatrix} 0 \\ * \\ * \end{pmatrix}$

where  $*$  entries are non-zero. Continuing in this way, we see that if  $M$  is  $\mathbb{F}$ -representable, it has an  $\mathbb{F}$ -representation of the form

$$\left[ \begin{array}{ccc|ccc} a & b & c & d & e & f & g \\ 1 & 0 & 0 & 0 & * & * & * \\ 0 & 1 & 0 & * & 0 & * & * \\ 0 & 0 & 1 & * & * & 0 & * \end{array} \right]$$

where  $*$ s are non-zero.

By rescaling rows and columns, there is an  $\mathbb{F}$ -representation



of the form:

$$\begin{bmatrix} a & b & c & d & e & f & g \\ 1 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 & z & 1 \\ 0 & 0 & 1 & x & y & 0 & 1 \end{bmatrix}$$

where  $x, y$  and  $z$  are non-zero.

Since  $\{a, d, g\}$  is dependent,  $\begin{vmatrix} 1 & 1 \\ x & 1 \end{vmatrix} = 0$ , so  $x=1$ .

Since  $\{b, e, g\}$  is dependent,  $\begin{vmatrix} 1 & 1 \\ y & 1 \end{vmatrix} = 0$ , so  $y=1$ .

Since  $\{c, f, g\}$  is dependent,  $\begin{vmatrix} 1 & 1 \\ z & 1 \end{vmatrix} = 0$ , so  $z=1$ .

When  $M = F_7$ ,  $\{d, e, f\}$  is dependent, so



$$\begin{aligned}
 \begin{vmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{vmatrix} &= - \begin{vmatrix} 1 & 1 \\ 1 & 0 \end{vmatrix} + \begin{vmatrix} 1 & 1 \\ 0 & 1 \end{vmatrix} \\
 &= 1 + 1 \\
 &= 0,
 \end{aligned}$$

so this method can only be represented over fields with characteristic 2 ( $\text{GF}(2)$  or  $\text{GF}(2^k)$ ) for  $k \geq 1$

After checking the other 3-element sets, we see that we have constructed a  $\mathbb{F}$ -representation for  $F_7$  provided  $\mathbb{F}$  has characteristic 2.

For the non-Fano method,

$$A = \left[ I \begin{vmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{vmatrix} \right] \mapsto \text{an } \mathbb{F}\text{-representation}$$

provided  $\mathbb{F}$  has characteristic not 2.

